

# CS356: More Networking

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# Firewalls: Ethernet Frames

preamble	SED	Destination	Source	tag	length	payload	CRC
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- Preamble (7 bytes) Alternating 0s and 1s; allows for synchronization
- Start of frame delimiter (1 byte) Indicates upcoming bits start frame
- Destination address (6 bytes) MAC address of destination
- Source address (6 bytes) MAC address of source
  - Tag (4 bytes) An optional tag
  - Length (2 bytes) Length of entire Ethernet frame
  - Payload (46–1500 bytes) Data carried by frame (e.g., IP packet)
- Cyclic redundancy check (4 bytes) Hash of addresses, length, and data

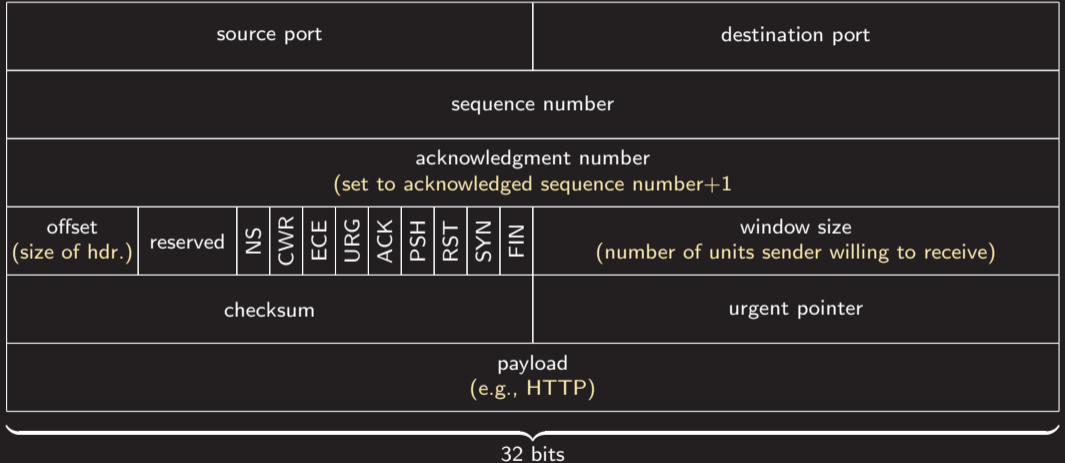
# Firewalls: Internet Protocol Packets



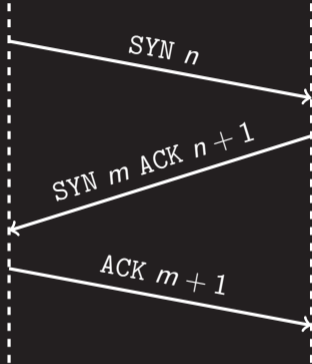
version 4 for IPv4	IHL (header len.)	DSCP (e.g., realtime)	ECN	total length (including header and payload)	
identification (collect fragments)			flags	fragment offset	
time to live (avoid circles)	protocol (e.g., 6 is TCP)		header checksum		
source IP address					
destination IP address					
options (if IHL>5)					
payload (e.g., TCP datagram)					

32 bits

# Firewalls: Transmission Control Protocol Segments



# More Networking: TCP Details



- ▶ Three-way handshake establishes initial sequence number, providing weak authenticator and preparing for reliability (`connect` system call).
- ▶ Subsequent segments contain SYN and ACK values (`recv/send/read/write`).
- ▶ Sender will at some point stop sending to allow ACKs to catch up (window size; transparent buffering).
- ▶ FIN flag indicates connection should close (`close/shutdown`).
- ▶ RST flag indicates something went wrong.

Wireshark demo: HTTPS (filter: `tcp.port==443`)

- ▶ Send SYN but never ACK.
- ▶ Guess of observe sequence number.
  - Sent FIN segments.
  - Sent RST segments.
  - Inject data segments.

Protection: SYN cookies (DJB, 1996)

$t$ : timestamp

$m$ : maximum segment size normally stored in SYN queue

$F$ : server-selected secret function

$s$ :  $F(l_1, P_1, l_2, P_2, t)$

ISN:  $t \bmod 32 \parallel m \parallel s$

- 1 Check freshness of  $t$
- 2 Recompute  $s$  and compare
- 3 Decode  $m$

<https://cr.yp.to/syncookies.html>

# More Networking: How Fast Can We Go?



source port	destination port
length	(optional) checksum

- ▶ DNS
- ▶ Real-time audio streams
- ▶ Real-time video streams

Wireshark demo: DNS (filter: dns)



# More Networking: Port Scanning

TCP (nmap's `-sT`) Complete connection, but slow and noisy

SYN (nmap's `-sS`) Send SYN watch for SYN-ACK

UDP (nmap's `-sU`) Send datagram, watch for ICMP port unreachable

FIN (nmap's `-sF`) Send spurious FIN; might traverse firewall; closed ports send RST, open ports ignore

See `nmap` man page.





# More Networking: UDP Programming

- ① `connect` has a different meaning: the protocol itself does not have the concept of a connection, so `connect` simply indicates the host to which subsequent `write/send` operations will transmit or from where `read/recv` will accept datagrams. Alternatively, skip `connect` and use `sendto/recvfrom`.
- ② `getattrinfo` will produce translations relevant for UDP if you set the hint parameter's `ai_socktype` field to `SOCK_DGRAM`.
- ③ We have sometimes relied on the end-of-file condition as a protocol terminator. With no connection, there is no end-of-file. Instead, the structure of the message must indicate when a protocol is complete. Common approaches include a special delimiter or a known length.

Complete `networkudp` in C on Aquinas.



Complete networkudp in C on Aquinas.

- ▶ Consolidate and reorganize network in C.
- ▶ Refine error handling.
- ▶ Provide clear usage message, e.g., "Usage: %s HOST PORT\n", argv[0].
- ▶ Set ai\_socktype field to SOCK\_DGRAM in hint parameter to getattrinfo.
- ▶ Rethink read/recv loop.



# More Networking: Assignments

Graded Homework Aquinas: udpscan

Reading Read 0x400–0x433

<https://www.flyn.org/courses/cs356/schedule>